

<<算法技术手册>>

图书基本信息

书名：<<算法技术手册>>

13位ISBN编号：9787564116323

10位ISBN编号：7564116323

出版时间：2009-4

出版时间：东南大学出版社

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页数：343

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## &lt;&lt;算法技术手册&gt;&gt;

## 内容概要

创造稳定的软件需要有效的算法，但是程序设计者们很少能在问题出现之前就想到。

《算法技术手册（影印版）》描述了现有的可以解决多种问题的算法，并且能够帮助你根据需求选择并实现正确的算法——只需要一定的数学知识即可理解并分析算法执行。

相对于理论来说，本书更注重实际运用，书中提供了多种程序语言中可用的有效代码解决方案，可轻而易举地适合一个特定的项目。

有了这本书，你可以：  
    解决特定编码问题或改进现有解决方案的执行；  
    迅速确定与需要解决的问题相关的算法，并判定为什么这样的算法是正确的；  
    探索C、C++、Java、Ruby中的算法解决方案，伴有实现诀窍；  
    了解一个算法预期的执行情况最佳的执行条件；  
    发现不同算法中相似设计产生的冲突；  
    学习先进的数据结构以改进算法效率。

有了《算法技术手册》，你可以学习如何改进算法的性能，这是软件应用成功的关键。

## 作者简介

George T. Heineman, Gary Pollice和Stanley Selkow均为 Worcester Polytechnic Institute (伍斯特理工学院) 计算机科学系的教授。

George是《Component-Based Software Engineering: Putting the Pieces Together》(Addison-Wesley) 的合编者, Gary则是《Head First Object-Oriented Analysis and Design》(O'Reilly) 的合著者。

## 书籍目录

Part 11. Algorithms Matter Understand the Problem Experiment if Necessary Algorithms to the Rescue Side Story The Moral of the Story References 2. The Mathematics of Algorithms Size of a Problem Instance Rate of Growth of Functions Analysis in the Best, Average, and Worst Cases Performance Families Mix of Operations Benchmark Operations One Final Point References 3. Patterns and Domains Patterns: A Communication Language Algorithm Pattern Format Pseudocode Pattern Format Design Format Empirical Evaluation Format Domains and Algorithms Floating-Point Computations Manual Memory Allocation Choosing a Programming Language References

Part 24. Sorting Algorithms Overview Insertion Sort Median Sort Quicksort Selection Sort Heap Sort Counting Sort Bucket Sort Criteria for Choosing a Sorting Algorithm References

5. Searching Overview Sequential Search Binary Search Hash-based Search Binary Tree Search 6. Graph Algorithms Overview Depth-First Search Breadth-First Search Single-Source Shortest Path All Pairs Shortest Path Minimum Spanning Tree Algorithms References

7. Path Finding in AI Overview Depth-First Search Breadth-First Search A Search Comparison Minimax NegMax AlphaBeta References

8. Network Flow Algorithms Overview Maximum Flow Bipartite Matching Reflections on Augmenting Paths Minimum Cost Flow Transshipment Transportation Assignment Linear Programming References

9. Computational Geometry Overview Convex Hull Scan Line Sweep Nearest Neighbor Queries Range Queries References

Part 310. When All Else Fails Variations on a Theme Approximation Algorithms Offline Algorithms Parallel Algorithms Randomized Algorithms Algorithms That Can Be Wrong, but with Diminishing Probability References

11. Epilogue Overview Principle: Know Your Data Principle: Decompose the Problem into Smaller Problems Principle: Choose the Right Data Structure Principle: Add Storage to Increase Performance Principle: If No Solution Is Evident, Construct a Search Principle: If No Solution Is Evident, Reduce Your Problem to Another Problem That Has a Solution Principle: Writing Algorithms Is Hard——Testing Algorithms Is Harder

Part 4 Appendix: Benchmarking Index

## 章节摘录

In the sortPointers function of Example 4 - 11, each element in the input is inserted into its associated bucket based upon the provided hash function; this takes linear, or  $O(n)$ , time. The elements in the buckets are not sorted, but because of the careful design of the hash function, we know that all elements in bucket  $b_j$  are smaller than the elements in bucket  $b_j$ , if  $i < j$ . As the values are extracted from the buckets and written back into the input array, INSERTION SORT is used when a bucket contains more than a single element. For BUCKET SORT to exhibit  $O(n)$  behavior, we must guarantee that the total time to sort each of these buckets is also  $O(n)$ . Let  $n_i$  define  $n_i$  to be the number of elements partitioned in bucket  $b_i$ . We can treat  $n_i$  as a random variable (using statistical theory). NOW consider the expected value. Each element in the input set has probability  $p = 1/n$  of being inserted into a given bucket because each of these elements is uniformly drawn from the range  $[0, 1)$ . Therefore,  $E[n_i] = n * p = n * (1/n) = 1$ . From this equation we can compute the expected value of  $n_i^2$ . This is critical because it is the factor that determines the COST of INSERTION SORT, which runs in a worst case of  $O(n^2)$ . We compute  $E[n_i^2] = (1 - 1/n) + 1 = (2 - 1/n)$ , which shows that  $E[n_i^2]$  is a constant. This means that when we sum up the COSTS of executing INSERTION SORT on all  $n$  buckets, the expected performance COST remains.

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### 媒体关注与评论

“作者汲取了大量鲜为人知的文献资料，这本不可或缺的指南巩固了理论与实际操作的完美平衡。通过它来理解算法变得更加轻松容易。”  
——Matthew Russell，高级技术总监，Digital Reasoning System；《Dojo：The Definitive Guide》的作者（OReilly）

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